Dynamic Replication and Partitioning

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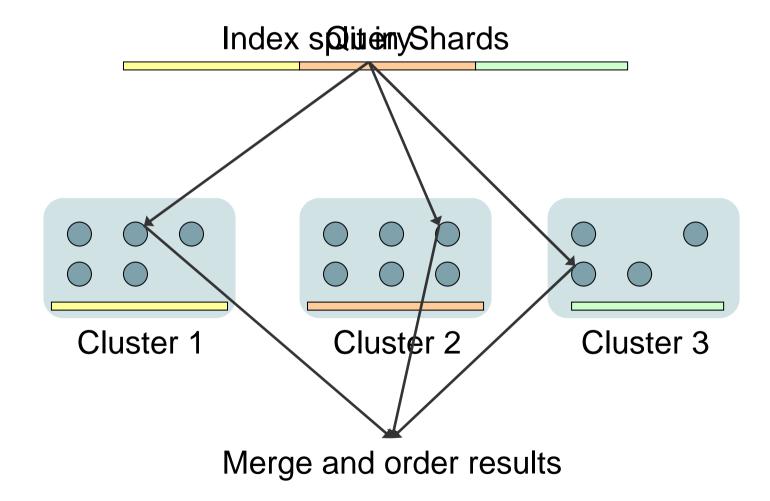
Joint work with Mark Handley, David S. Rosenblum

Motivation: Web Search

- Search engines
 - Create an index of the web
 - Queries consult the index to find relevant documents
 - The documents are then ordered (e.g. Page Rank)
- The index is huge: a few TB
 - Must be partitioned to fit into memory
 - Must be replicated to increase query throughput and system availability

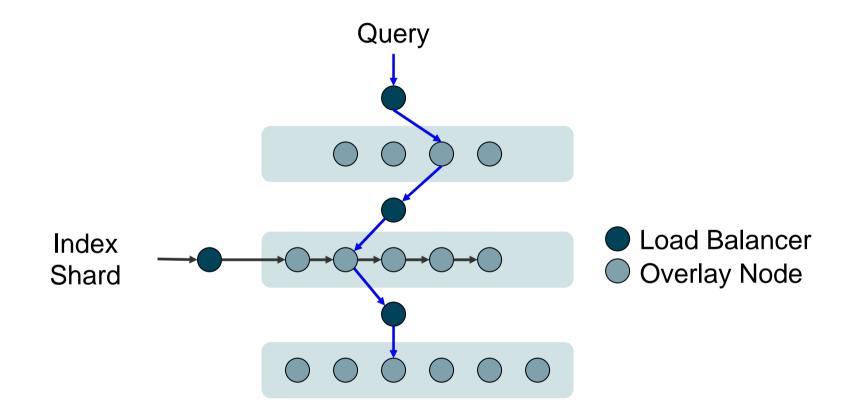


Google Web Search (Barroso et. al)





Big Picture: Distributed Rendez-Vous



Average Replication Level R=5 Hop Count H=3

Distributed Rendez Vous is important

- Many other applications use it
 - Online Filtering
 - Distributed databases
- Combines replication and partitioning
 - Increasing replication (R) increases availability, but has high cost for storing the index
 - Increasing the forwarding hops (H) creates high bandwidth cost for transient objects
 - − Tradeoff: $R \cdot H \ge #$ nodes

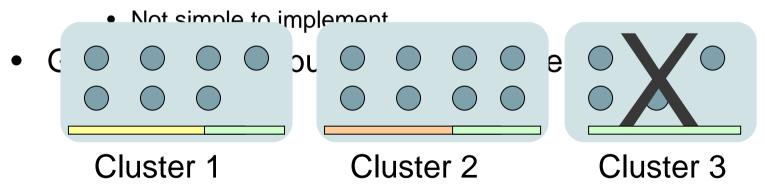
The Problem

- Who chooses the number of clusters? Depends on:
 - Frequencies and sizes of index and queries
 - Bandwidth constraints
 - Memory constraints
 - Number of nodes
- R varies with time!

How can we adjust the *Replication Rate* in distributed rendez-vous?

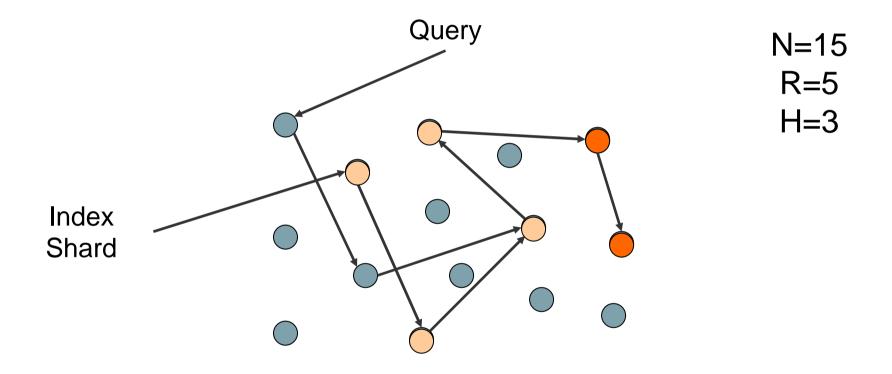
Obvious approach

- Google architecture
 - Replication tied to network structure
 - Increase replication level
 - Destroy cluster, add the nodes to the other clusters
 - Issues
 - Temporarily reduces the capacity of the network





A randomized implementation



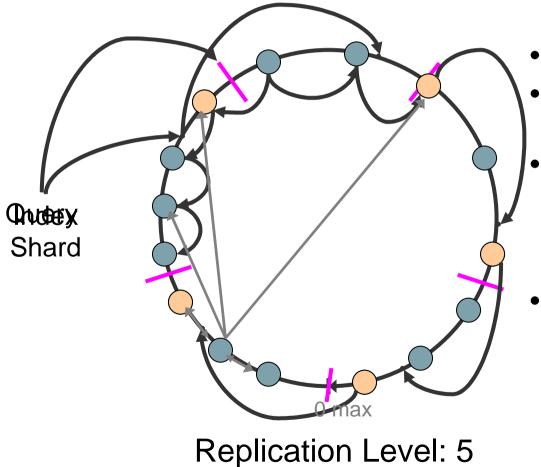
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Our solution: ROAR

• Rendez-Vous On A Ring

- Similar in spirit to Random
- But with deterministic properties
- Does not tie network structure to replication level

ROAR Overview



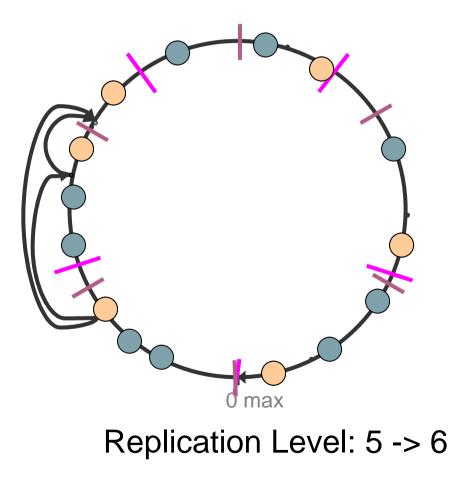
- Nodes on a Chord ring
- ID space virtually split in R intervals
- Replicate
 - Hash and store
 - Forward to equivalent node in next interval
- Route
 - Uniformly choose interval and direction
 - Route to all nodes in that interval

ROAR Analysis

- Equal spacing is important
 - When R increases, it ensures that no 2 replicas are in the same interval
 - Stable state: if R is constant enough time, equivalent nodes have equivalent content
 - Useful for fault tolerance
 - When R changes:
 - Stability is maintained if R is doubled of halved
 - Otherwise, not stable: wait for objects to expire



Increasing Replication



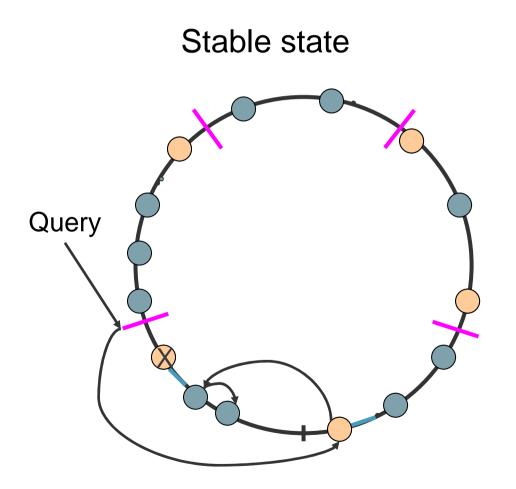
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Increasing Replication (2)

- **Observation**. When replication level is R, we can route at any level R'≤R.
- ROAR can route while changing replication levels
 - Wait until all nodes in interval reach new replication
 - Begin routing at new replication level
- When is the new replication level reached?
 - Compute persistent object count at replication level R and R+1
 - When approximately equal, safe to switch to new routing.
 - Count is piggybacked on queries very small cost

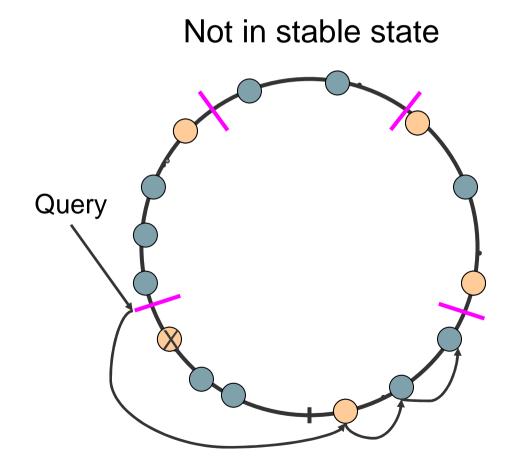


Fault Tolerance



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Fault Tolerance



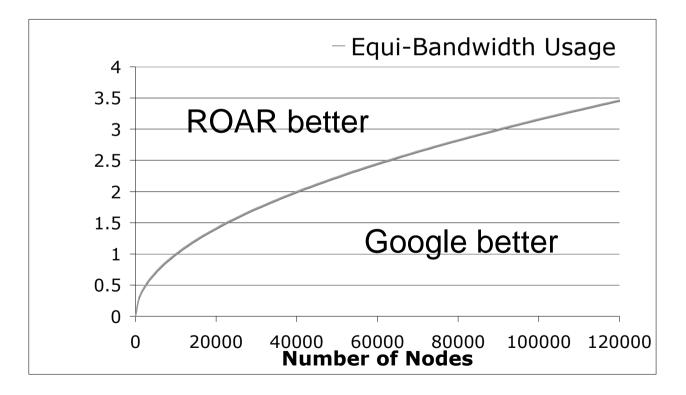
Comparison

- Bandwidth scarce system
 - $R = O(\sqrt{N})$
 - I = # total size of index

	Google	Random	ROAR
RV Guaranteed?	Yes	35% miss probability	Yes
RV Redundant?	No	25% redundant RV probability	No
Bw for R = R+1	~2.1	Ι	
Bw Cost on Node Failure	1	O(I·R/N)	O(I·R/N) or 1

Comparison (2)

- 1% permanent failures per year
 - Commercial data: 5% failures in 1st year
 - Transient failures tolerated with stable state



Summary

- Distributed rendez-vous is an important problem in distributed computing
 - Changing R is a requirement for optimal solutions
- ROAR simple algorithm
 - Distributed in spirit
 - No need for external load balancing
 - Can run on deployed structured overlays
 - Achieves reconfiguration without changing network structure
 - In stable state as good as Google
 - When reconfigurations are often, does better

References

• Web Search for a Planet: the Google Cluster Architecture - Barroso et. al