

# A Continuous Dcpo Representation of Regular Formal Topologies

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Domains IX

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## Background/Motivation

- ▶ Formal topology includes a constructive theory of domains, in the sense that the formal space  $Pt(\mathcal{X})$  of a *Scott (or unary) formal topology*  $\mathcal{X}$  is a Scott domain when ordered by inclusion, and every Scott domain arises in this way (G. Sambin, 1987).

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- ▶ In fact, the formal space  $Pt(\mathcal{X})$  of  $\mathcal{X}$  can be embedded in a Scott domain  $D = (Pt(\mathcal{X}_S), \subseteq)$ , where  $\mathcal{X}_S$  is the Scott compactification of  $\mathcal{X}$  (so  $\mathcal{X}_S$  is Scott) (G. Sambin 1987). Hence, the Scott compactification gives a domain representation of the topological space  $Pt(\mathcal{X})$ , giving a connection between the two approaches of representing topological spaces (E. Palmgren, 2007).

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- ▶ This representation comes with an almost automatic lifting of morphisms: If  $F : \mathcal{X} \rightarrow \mathcal{Y}$  continuous, then  $F_S =_{\text{def}} F : \mathcal{X}_S \rightarrow \mathcal{Y}_S$  is continuous and the induced continuous function  $Pt(F_S) : Pt(\mathcal{X}_S) \rightarrow Pt(\mathcal{Y}_S)$  satisfies  $Pt(F_S) \upharpoonright_{Pt(\mathcal{X})} = Pt(F)$  (E. Palmgren, 2007).

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- ▶ There is another result saying that every *stable locally Scott* formal topology  $\mathcal{X}$  give rise to a continuous dcpo  $(Pt(\mathcal{X}), \subseteq)$ , and that every continuous dcpo arises in this way (S. Negri 2002). There is a nice extension of the reals in this class of formal topologies, called *the partial reals* in (S. Negri, 2002).

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- ▶ The aim now is to describe a continuous domain representation of formal topologies as continuous dcpo's, in such a way that the formal reals are represented in the partial reals. To make this work we have to restrict the class of formal topologies.

# Formal Topology

## Definition

A *Formal Topology* (with positivity predicate) is a tuple  $\mathcal{X} = (X, \leq, \triangleleft, \text{Pos})$  where  $(X, \leq)$  is a pre-ordered set (the set of *basic opens*) and  $\triangleleft$  is a relation between basic opens  $a \in X$  and subsets  $U \subseteq X$  satisfying

- (Ref) If  $a \in U$ , then  $a \triangleleft U$ ,
- (Tra) If  $a \triangleleft U$  and  $U \triangleleft V$ , then  $a \triangleleft V$ ,
- (Ext) If  $a \leq b$ , then  $a \triangleleft \{b\}$ ,
- (Loc) If  $a \triangleleft U$  and  $a \triangleleft V$ , then  $a \triangleleft U_{\leq} \cap V_{\leq}$ .



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- (Loc) If  $a \triangleleft U$  and  $a \triangleleft V$ , then  $a \triangleleft U_{\leq} \cap V_{\leq}$ .

Here  $U \triangleleft V$  just means that  $u \triangleleft V$  for all  $u \in U$  and  $U_{\leq} =_{\text{def}} \{x \in X : (\exists u \in U)x \leq u\}$ .

If  $a \triangleleft U$  we say that  $U$  covers  $a$ , and we call  $\triangleleft$  a cover relation.

## Formal Topology

Furthermore, Pos is a predicate on  $X$  satisfying

(Mon) If  $\text{Pos}(a)$  and  $a \triangleleft U$ , then there is  $b \in U$  such that  $\text{Pos}(b)$ ,

(Pos) For all  $a \in X$ ,  $a \triangleleft \{a\}^+$ ,

where  $U^+ =_{\text{def}} \{x \in X : x \in U \ \& \ \text{Pos}(x)\}$ .

# Formal Topology

## Definition

A (*formal*) *point* in a formal topology  $\mathcal{X}$  is a subset  $\alpha \subseteq X$  satisfying

- (i) there is  $a \in \alpha$ ,
- (ii)  $a, b \in \alpha$  iff there is  $c \in \alpha$  such that  $c \leq a, b$ ,
- (iii) if  $a \in \alpha$  and  $a \triangleleft U$ , then there is  $b \in U \cap \alpha$ ,
- (iv) if  $a \in \alpha$  then  $\text{Pos}(a)$ .

The clause (iv) can actually be derived using (iii) and (Pos) ( $a \triangleleft a^+$  for all  $a \in X$ ), but we will need it later.

## Formal Topology

The collection of points in a formal topology  $\mathcal{X}$  is denoted  $Pt(\mathcal{X})$ , and the collection  $\text{ext}_{\mathcal{X}}(U)$  of points associated with a set  $U$  of basic opens is given by

$$\text{ext}_{\mathcal{X}}(U) =_{\text{def}} \{\alpha \in Pt(\mathcal{X}) : (\exists a \in U)a \in \alpha\}.$$

The collections  $\text{ext}_{\mathcal{X}}(U)$ ,  $U \subseteq X$ , form an ordinary topology on  $Pt(\mathcal{X})$ .

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## Theorem

*If  $\mathcal{X}$  is a formal topology, then  $(Pt(\mathcal{X}), \subseteq)$  is a directed complete partial order (dcpo).*

# CLSP Formal Topologies

## Definition

We say that a pre-order  $(X, \leq_{\mathcal{X}})$  is a *consistently complete lower semi-lattice pre-order* (CLSP), if  $X$  has a largest element  $\top$  and if for every consistent pair  $a, b \in X$ , i.e. which has a lower bound  $c \leq a$  and  $c \leq b$ , there is an element  $a \wedge b \in X$  satisfying

- (a)  $a \wedge b \leq a, b$ ,
- (b)  $c \leq a, b$  implies  $c \leq a \wedge b$ .

Then  $\mathcal{X} = (X, \leq, \triangleleft, \text{Pos})$  is called a CLSP formal topology whenever  $(X, \leq)$  is a CLSP.

## Scott compactification

Given a formal topology  $\mathcal{X} = (X, \leq, \triangleleft, \text{Pos})$  we define a new structure  $\mathcal{X}_S$ , as follows: First we define a new preorder  $\leq_S$  on  $X$  by

$$a \leq_S b \iff_{\text{def}} a \triangleleft b,$$

then we define a new relation

$$a \triangleleft_S U \iff_{\text{def}} \text{Pos}(a) \rightarrow \exists b \in U : a \leq_S b,$$

and define  $\mathcal{X}_S = (X, \leq_S, \triangleleft_S, \text{Pos})$ . One can show that  $\mathcal{X}_S$  is a formal topology, and we call it the *Scott compactification* of  $\mathcal{X}$  (S. Negri, 2002).

## Scott compactification

If  $\mathcal{X}$  is a formal topology, such that  $\mathcal{X}_S$  is a CLSP formal topology, then

- ▶  $(Pt(\mathcal{X}_S), \subseteq)$  is a Scott domain,
- ▶  $Pt(\mathcal{X}) \subseteq Pt(\mathcal{X}_S)$  and
- ▶ the point topology on  $Pt(\mathcal{X}_S)$  is the Scott Topology.



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- ▶ the point topology on  $Pt(\mathcal{X}_S)$  is the Scott Topology.

This can then be stated in terms of a (homeomorphic) domain representation  $(D, D^R, \varphi)$ , with

- ▶ domain  $D = (Pt(\mathcal{X}_S), \subseteq)$ ,
- ▶ representing elements  $D^R = Pt(\mathcal{X})$  and
- ▶ representation map  $\varphi = id_{Pt(\mathcal{X})} : D^R \rightarrow Pt(\mathcal{X})$ .

# Regular Formal Topologies

Given a formal topology  $\mathcal{X}$  we can define a relation  $\prec$  on  $X$  by

$$a \prec b \iff_{\text{def}} X \triangleleft a^\perp \cup \{b\},$$

where  $a^\perp =_{\text{def}} \{c \in X : c \leq \cap a \leq \triangleleft \emptyset\}$  (the *open complement* of  $a$ ).

If  $a \prec b$  we say that  $a$  is *well covered* by  $b$ .

# Regular Formal Topologies

## Definition

A formal topology  $\mathcal{X}$  is said to be *regular* if

$$b \triangleleft \{x \in X : x \prec b\},$$

for all  $b \in X$ .

Moreover, we will say that  $\prec$  is *dense* on  $X$  if

$$a \prec b \implies (\exists c \in X) a \prec c \prec b.$$

# Regular Formal Topologies

Given a regular CLSP formal topology  $\mathcal{X}$  with  $\prec$  dense on  $X$ , we can define a new topology  $\mathcal{X}_R$  as follows: Let

$$a \leq_R b \iff_{\text{def}} a \leq_S b,$$

i.e.  $a \leq_R b$  iff  $a \triangleleft b$ , and

$$a \triangleleft_R U \iff_{\text{def}} \{c \in X : c \prec a\} \triangleleft_S U,$$

i.e.  $a \triangleleft_R U \iff (\forall c \prec a)(\text{Pos}(c) \rightarrow (\exists b \in U)(c \triangleleft b))$ .

We set  $\mathcal{X}_R = (X, \leq_R, \triangleleft_R, \text{Pos})$ .

# Regular Formal Topologies

## Proposition

*If  $\mathcal{X}$  is regular, CLSP and  $\prec$  is dense on  $X$ , then  $\mathcal{X}_R$  is a formal topology.*

The topology  $\mathcal{X}_R$  is not regular but it (trivially) satisfies

$$b \triangleleft_R \{a \in X \mid a \prec b\}$$

for all  $b \in X$ .

Note: the relation  $\prec$  is still defined via the original cover  $\triangleleft$ .

## Points of $\mathcal{X}_R$

### Lemma

The points of  $\mathcal{X}_R$  are precisely those subsets  $\alpha \subseteq X$  satisfying

- (i) There is  $a \in \alpha$ ,
- (ii)  $a, b \in \alpha$  iff there is  $c \in \alpha$  such that  $c \leq_R a, b$ ,
- (iii) If  $a \in \alpha$ , then there is  $c \in \alpha$  with  $c \prec a$ ,
- (iv)  $\text{Pos}(a)$  for all  $a \in \alpha$ .

For any  $a \in X$  with  $\text{Pos}(a)$  we have

$$\uparrow a =_{\text{def}} \{x \in X : a \prec x\} \in \text{Pt}(\mathcal{X}_R)$$

# Regular Formal Topologies and Continuous Dcpo's

## Theorem

If  $\mathcal{X}$  is regular, CLSP and  $\prec$  is dense on  $X$ , then

- (a)  $Pt(\mathcal{X}_R)$  is a continuous dcpo with a base given by elements  $\uparrow a$  with  $Pos(a)$ ,
- (b)  $Pt(\mathcal{X}) \subseteq Pt(\mathcal{X}_R)$ ,
- (c)  $Pt(\mathcal{X}) \cap ext_{\mathcal{X}_R}(U) = ext_{\mathcal{X}}(U)$ .

The way below relation  $\ll$  on  $Pt(\mathcal{X}_R)$  has a simple characterization:

$$\alpha \ll \beta \iff (\exists b \in \beta)(\alpha \subseteq \uparrow b \subseteq \beta).$$

## Regular Formal Topologies and Continuous Dcpo's

The point topology on  $Pt(\mathcal{X}_R)$  is precisely the Scott topology, so we can again state the above result as a domain representation  $(D, D^R, \varphi)$ , where

- ▶  $D = Pt(\mathcal{X}_R)$ ,
- ▶  $D^R = Pt(\mathcal{X})$  and
- ▶  $\varphi = id_{Pt(\mathcal{X})} : D^R \rightarrow Pt(\mathcal{X})$ .

As in the case of the Scott compactification, we get a lifting of continuous morphism (although not as automatic).



# Continuous Morphisms

## Definition

A *continuous morphism*  $F : \mathcal{X} \rightarrow \mathcal{Y}$  between formal topologies is a relation  $F \subseteq X \times Y$  satisfying

- (A1)  $aFb, b \triangleleft_{\mathcal{Y}} V \implies a \triangleleft_{\mathcal{X}} F^{-1}V,$
- (A2)  $a \triangleleft_{\mathcal{X}} U, xFb \text{ for all } x \in U \implies aFb$
- (A3)  $X \triangleleft_{\mathcal{X}} F^{-1}Y,$
- (A4)  $aFb, aFc \implies a \triangleleft_{\mathcal{X}} F^{-1}(b_{\leq \mathcal{Y}} \cap c_{\leq \mathcal{Y}})$

Here  $F^{-1}Z =_{\text{def}} \{x \in X : (\exists z \in Z)xFz\}.$

# Continuous Morphisms

Each continuous morphism  $F : \mathcal{X} \rightarrow \mathcal{Y}$  induces a continuous point function  $Pt(F) : Pt(\mathcal{X}) \rightarrow Pt(\mathcal{Y})$  defined by

$$Pt(F)(\alpha) = \{b \in Y \mid \exists a \in \alpha : aFb\}.$$

A continuous morphism  $F$  then satisfies

$$aFb \implies Pt(F)(\text{ext}_{\mathcal{X}}(a)) \subseteq \text{ext}_{\mathcal{Y}}(b),$$

for all  $a \in X$ ,  $b \in Y$ .

# Continuous Morphisms

Moreover, point functions are monotone

$$\alpha \sqsubseteq \beta \implies Pt(F)(\alpha) \sqsubseteq Pt(F)(\beta),$$

and preserve directed suprema

$$\alpha_i \ (i \in I) \text{ directed} \implies Pt(F)\left(\bigcup_{i \in I} \alpha_i\right) = \bigcup_{i \in I} Pt(F)(\alpha_i).$$

So  $Pt(F)$  is Scott continuous between dcpos  
 $(Pt(\mathcal{X}), \sqsubseteq) \rightarrow (Pt(\mathcal{Y}), \sqsubseteq)$ .

## Lifting Continuous Morphisms

Every continuous morphism  $F : \mathcal{X} \rightarrow \mathcal{Y}$  between regular CLSP formal topologies, with dense well covered relations, induces a new relation  $F_R \subseteq X \times Y$  defined by

$$aF_Rb \iff_{\text{def}} \{x \in X : x \prec a\}^+ \subseteq F^{-1}\{y \in Y : y \prec b\}.$$

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## Theorem

*The relation  $F_R$  is a continuous morphism between  $\mathcal{X}_R$  and  $\mathcal{Y}_R$ , and hence  $Pt(F_R) : Pt(\mathcal{X}_R) \rightarrow Pt(\mathcal{Y}_R)$  is Scott continuous.*

*Moreover, if  $\alpha \in Pt(\mathcal{X}) \subseteq Pt(\mathcal{X}_R)$  then*

$$Pt(F_R)(\alpha) = Pt(F)(\alpha).$$

## Formal Reals $\mathcal{R}$

The basic opens of  $\mathcal{R}$  are all pairs

$$(p, q) \in (\mathbb{Q} \cup \{-\infty\}) \times (\mathbb{Q} \cup \{+\infty\}),$$

where  $p < q$ . The pre-order is given by inclusion (as intervals), i.e.

$$(p, q) \leq (r, s) \iff_{\text{def}} r \leq p < q \leq s.$$

Then  $\mathcal{R}$  is a CLSP formal topology with top element  $(-\infty, +\infty)$  and  $\wedge$  given by interval intersection

$$(p, q) \wedge (r, s) =_{\text{def}} (\min(p, r), \max(q, s)).$$

## Formal Reals $\mathcal{R}$

The cover  $\triangleleft$  of  $\mathcal{R}$  is generated by

(G1)  $(p, q) \triangleleft \{(p', q') : p < p' < q' < q\}$ , all  $(p, q)$ ,

(G2)  $(p, q) \triangleleft \{(p, s), (r, q)\}$ , all  $(p, q)$  and  $p \leq r < s \leq q$ .

The positivity predicate  $\text{Pos}$  is trivial since we only consider  $(p, q)$  with  $p < q$ .

The space  $Pt(\mathcal{R})$  is homeomorphic to the ordinary real numbers (S. Negri, 1996).

## Formal Reals $\mathcal{R}$

The well covered relation  $\prec$  on basic opens is given by

$$(p, q) \prec (r, s) \iff (r < p \vee r = -\infty) \ \& \ (q < s \vee s = +\infty).$$

Hence,  $\mathcal{R}$  is a regular CLSP formal topology and  $\prec$  is dense.

Applying the previous results we get an extension  $Pt(\mathcal{R}_R) \supseteq Pt(\mathcal{R})$  which is a continuous dcpo when ordered by inclusion.

What are the points  $Pt(\mathcal{R}_R)$  of  $\mathcal{R}_R$ ?



## The dcpo $Pt(\mathcal{R}_R)$

One can show that  $Pt(\mathcal{R}_R)$  is (essentially) the structure called *the partial reals* in (S. Negri, 2002). These are the points of the formal topology  $\mathcal{R}_p$  you get by removing (G2) from the generation of the cover of  $\mathcal{R}$ .

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Another description is given by the generalized reals in (F. Richman, 1998). Richman's generalized reals are given as pairs,  $(L, U)$ , of disjoint open subsets,  $L, U \subseteq \mathbb{Q}$  satisfying  $p < q \in L \implies p \in L$  and  $p > q \in U \implies p \in U$ . These are ordered by inclusion:

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$$(L, U) \leq (L', U') \iff_{\text{def}} L \subseteq L' \ \& \ U' \subseteq U.$$

There is a bijective correspondence between such cuts and elements of  $Pt(\mathcal{R}_R)$ . The above ordering of cuts induces a partial order on  $Pt(\mathcal{R}_R)$  that coincides with usual partial order of formal reals.

## The dcpo $Pt(\mathcal{R}_R)$

Some examples of points:

$$-\infty =_{\text{def}} \{(-\infty, q) : q \in \mathbb{Q}\} \text{ and } +\infty =_{\text{def}} \{(p, +\infty) : p \in \mathbb{Q}\}.$$

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If we for every  $\alpha \leq \beta \in Pt(\mathcal{R})$  define

$$I[\alpha, \beta] =_{\text{def}} \{(p, q) : (p, +\infty) \in \alpha \ \& \ (-\infty, q) \in \beta\},$$

then  $I[\alpha, \beta] \in Pt(\mathcal{R}_R)$ .

Points  $I[p, q]$  with  $p < q \in \mathbb{Q} \cup \{-\infty, +\infty\}$  make up the base of  $Pt(\mathcal{R}_R)$ .

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Distances between points and sets in a metric space (F. Richman, 1998).

## Continuous Functions on $Pt(\mathcal{R})$

We know that every continuous morphism  $F : \mathcal{R} \rightarrow \mathcal{R}$  lift to a Scott continuous function  $Pt(F_R) : Pt(\mathcal{R}_R) \rightarrow Pt(\mathcal{R}_R)$  satisfying  $Pt(F_R)(\alpha) = Pt(F)(\alpha)$  for all  $\alpha \in Pt(\mathcal{R})$ .

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Now, every continuous function  $f : \mathbb{R} \rightarrow \mathbb{R}$  (in the sense of Bishop (Bishop, Bridges, 1985)) can be represented by a continuous morphism  $A_f : \mathcal{R} \rightarrow \mathcal{R}$  (Palmgren, 2004). Hence we can lift every such function to the dcpo  $Pt(\mathcal{R}_R)$  and apply it to nonstandard elements.



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### Theorem

*If  $f : \mathbb{R} \rightarrow \mathbb{R}$  is continuous and  $a \leq b \in \mathbb{R}$ , then  $i = \inf_{x \in [a, b]} f(x)$  and  $s = \sup_{x \in [a, b]} f(x)$  exist, and*

$$f_R(I[a, b]) = I[i, s].$$

## Arithmetic in $Pt(\mathcal{R}_R)$

Using this one can show that arithmetic on  $Pt(\mathcal{R}_R)$  coincides with interval arithmetic on interval points: for  $a \leq b, c \leq d \in \mathbb{R}$ , we have

$$I[a, b] +_R I[c, d] = I[a + c, b + d]$$

and

$$I[a, b] \times_R I[c, d] = I[\min(ac, ad, bc, bd), \max(ac, ad, bc, bd)].$$

## Domain Theoretic Derivative

In (A. Edalat & A. Lieutier, 2001) the authors introduced an operation  $D$  on functions  $f : \mathbb{I}\mathbb{R} \rightarrow \mathbb{I}\mathbb{R}$ , where  $\mathbb{I}\mathbb{R}$  consists of all closed real intervals together with  $\mathbb{R}$  ordered under reverse inclusion, satisfying

- $D(f)$  continuous for every continuous  $f : \mathbb{I}\mathbb{R} \rightarrow \mathbb{I}\mathbb{R}$ ,
- $D(\mathbf{I}f) = \mathbf{I}(\frac{df}{dx})$ , for all  $\mathcal{C}^1$  functions  $f : \mathbb{R} \rightarrow \mathbb{R}$ .

Here  $\mathbf{I}(f)$  is the lifting of  $f$  given by the direct image under  $f$ .






## Domain Theoretic Derivative

In our setting (i.e. replacing  $\mathbb{R}$  by  $Pt(\mathcal{R}_R)$ ) we can show that  $D(f_R)$  is well defined when  $f : \mathbb{R} \rightarrow \mathbb{R}$  is differentiable (again in the sense of Bishop) and






$$D(f_R) = \left( \frac{df}{dx} \right)_R.$$

Thank You!

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